



## Slack time evaluation with RTSJ

Damien Masson, Serge Midonnet

### ► To cite this version:

Damien Masson, Serge Midonnet. Slack time evaluation with RTSJ. SAC 2008 Poster Session, Mar 2008, Brazil. pp.322–323. hal-00620348

**HAL Id: hal-00620348**

**<https://hal.science/hal-00620348>**

Submitted on 30 Sep 2011

**HAL** is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

# Slack Time Evaluation with RTSJ

Damien Masson and Serge Midonnet  
Université Paris-Est  
Laboratoire d'informatique de l'institut Gaspard-Monge  
UMR 8049 IGM-LabInfo  
77454 Marne-la-Vallée Cedex 2, France  
{masson, midonnet}@univ-mlv.fr

## ABSTRACT

We address in this paper the problem of jointly scheduling hard periodic tasks and soft aperiodic events using the Real-Time Specification for Java (RTSJ). We present the programming constraints of RTSJ and propose slack time evaluation and utilization algorithms which take these constraints into account. We evaluate these algorithms and compare their performances with the Background Scheduling (BS) through simulations.

## 1. INTRODUCTION

The need for more flexibility in real-time systems leads the community to address the problem of jointly scheduling hard periodic tasks and soft aperiodic events.

The easiest way to achieve this is to schedule all non-periodic tasks at a lower priority (assuming that the tasks are scheduled using a preemptive fixed priority policy). This policy is known as the *Background Servicing (BS)*.

In order to minimize the aperiodic tasks response times whilst guaranteeing the feasibility of the periodic tasks, the periodic task servers were introduced. A server is a periodic task, for which classical response time determination methods are applicable (with or without modifications). The server is in charge of servicing the non-periodic traffic with a limited capacity.

The *Static Slack Stealer*, an algorithm to compute the exact slack available at time  $t$  to execute aperiodic tasks at the highest priority, is proposed in [1]. The optimality of this approach for maximizing the aperiodic tasks response times is demonstrated. Unfortunately, the time and memory complexities does not permit to use it.

We propose in this paper to start from RTSJ constraints to set up a slack time approximation algorithm. Then we discuss the slack time utilization issues with RTSJ.

We review in Section 2 the available mechanisms with RTSJ. Section 3 describes our approximate slack evaluation algorithm. Section 4 explicits RTSJ constraints to design slack time evaluation and utilization algorithms. Section

5 presents our simulations methodology and results. We conclude in Section 6.

## 2. APERIODIC TASKS AND RTSJ

The RTSJ proposes two classes `AsyncEvent` and `AsyncEventHandler` to model respectively an asynchronous event and its handler(s). The only way to include an handler in the feasibility process is to treat it as an independent task, and that implies to know at least its worst-case occurring frequency.

The RTSJ does not support any particular task server policy. We propose a solution in [2].

## 3. AVAILABLE SLACK EVALUATION

Our goal is to construct a slack time evaluation implementable with a minimal RTSJ compliant virtual machine.

The only instants in the system life time where we can easily operate are each periodic task instances begin and end.

We decompose the maximum available slack per task ( $S_{i,t}^{max}$ ) according to Equation 1 where  $W_{i,t}^{max}$  is the maximum possible work at priority  $i$  regardless of lower priority processes, and  $c_i(t)$  is the effective hard real-time work we have to process at the instant  $t$  at the priority  $i$ .

$$S_{i,t}^{max} = W_{i,t}^{max} - c_i(t) \quad (1)$$

Then the available slack at time  $t$  at the highest priority is the minimum of the  $S_{i,t}^{max}$ .

### 3.1 Data initialization

Under hypothesis of a synchronous activation of all periodic hard real-time tasks at  $t_0$ , we have:

$$\begin{cases} W_{1,t_0}^{max} &= d_1 \\ \forall i, W_{i,t_0}^{max} &= d_i - \sum_{\forall k \in HP(i)} \left\lceil \frac{T_i}{T_k} \right\rceil c_k \\ \forall i, c_i(t_0) &= c_i \end{cases} \quad (2)$$

where  $c_i$ ,  $d_i$  and  $T_i$  are respectively the cost, the deadline and the period of the periodic task  $\tau_i$ . The data initialization has a time complexity in  $O(n^2)$ , with  $n$  the periodic tasks number.

### 3.2 Dynamic operations

The data are function of time, so they potentially evolve at each clock tick. Between two dates  $t_1$  and  $t_2$ ,  $W_i$  is reduced by  $t_2 - t_1$  for any task and  $c_i(t)$  is reduced for any task executed during the interval. Eventually, these values can

Permission to make digital or hard copies of all or part of this work for personal or classroom use is granted without fee provided that copies are not made or distributed for profit or commercial advantage and that copies bear this notice and the full citation on the first page. To copy otherwise, to republish, to post on servers or to redistribute to lists, requires prior specific permission and/or a fee.

SAC'08 March 16-20, 2008, Fortaleza, Ceará, Brazil

Copyright 2008 ACM 978-1-59593-753-7/08/0003 ...\$5.00.

be increased for the tasks which complete between  $t_1$  and  $t_2$ . However, we only need to update them at each periodic task instance begin and end and when an aperiodic event occurs.

1. **End of periodic task  $\tau_i$ .** Let  $dt$  be the elapsed time since the last update of our data. Then,  $W_i$  is reduced by the elapsed time since the last update and increased by  $T_i$  minus the interference of tasks with higher priority. This interference, given by equation 3, has already been computed during the initialization phase.

$$I_i = \sum_{\forall k \in HP(i)} \left\lceil \frac{T_i}{T_k} \right\rceil c_k \quad (3)$$

$I_i$  is not the real interference at instant  $t$ , but the maximal interference. This bound is reached in the worst case scenario: the synchronous activation of all the tasks.

Moreover, for all tasks,  $W_k$  is reduced by the elapsed time, and for all tasks with a lower priority than  $\tau_i$ ,  $W_k$  is increased by  $c_i$ . Finally,  $c_i(t)$  is reset to  $c_i$ .

$$\begin{cases} \forall j, W_{j,t} = W_{j,t'} - dt \\ W_{i,t}^{max} \geq W_{i,t} = \max(W_{i,t'}, 0) + T_i - I_i \\ \forall j > i, W_{j,t} = W_{j,t'} + c_i \\ c_i(t) = c_i \end{cases} \quad (4)$$

2. **Beginning of periodic task  $\tau_i$ .**

The maximum available process time is reduced by  $dt$  at all priority level. Let  $j$  be the previous activity priority level. If  $j$  is in the hard real-time priority range (noted *hrtp*), then  $c_j(t)$  is reduced by  $dt$ .

$$\begin{cases} \forall k, W_{k,t} = W_{k,t'} - dt \\ \text{if } j \in \text{hrtp}, c_j(t) = c_j(t') - dt \\ \text{if } W_{i,t} < c_i(t), W_{i,t} = c_i(t) \end{cases} \quad (5)$$

3. **Soft real-time aperiodic task  $\tau_\alpha$  arrival.**

As for the previous cases, the maximum available process time is reduced by  $dt$  at all priority level as  $W_j$  if the previous activity priority,  $j$ , is in the hard real-time priority range. Then, the available slack at the highest priority is computed.

$$\begin{cases} \forall k, W_{k,t} = W_{k,t'} - dt \\ \text{if } j \in \text{hrtp}, c_j(t) = c_j(t') - dt \\ S_{i,t}^{max} \geq S_{i,t} = \min_{i \in \text{hrpr}} W_{i,t} - c_i(t) \end{cases} \quad (6)$$

## 4. RTSJ CONSTRAINTS

### 4.1 Slack approximation

In order to act at the beginning and at the end of the periodic tasks instances, we overload the `waitForNextPeriod()` method of `RealTimeThread`. This method is called at the end of each periodic job and is blocking until the next job execution.

The slack stealer may be implemented as an `AsyncEventHandler` bound to a special `AsyncEvent` and with the maximal available priority.

The soft aperiodic events can be implemented with several `AsyncEvent` all bound to the `AsyncEventHandler` representing the slack stealer. In this way, we can call the slack stealer every time a soft aperiodic event is released and before and after each instance of hard periodic tasks.

## 4.2 Slack Utilization

With Java, it is impossible to resume a previously stopped thread. The RTSJ offers a mechanism to asynchronously interrupt a RT-thread, but it does not allow to resume it later. Moreover, in most implementation of RTJVM, the dynamic change of the priority is not supported.

Due to these limitations, we cannot begin the execution of a soft task at the highest priority if there is not enough slack to permit its completion.

When there is available slack, we schedule the execution of the first soft pending task with a cost lower or equal than the available slack at the highest priority. When there is no slack and no pending periodic task, we schedule in BS the first pending aperiodic task.

## 5. SIMULATIONS

We simulate the same real-time systems with the three following aperiodic policies: a BS, a slack stealing with an exact computation of the slack and one with our approximation of the slack.

### 5.1 Methodology

We measure the mean response time of soft tasks with different aperiodic and periodic loads.

First, we generate groups of ten periodic task sets with utilization levels of 30, 50, 70 and 90%. The results presented in this section are based on averages over a group of ten task sets. Then, we generate groups of ten aperiodic task sets with a range of utilization levels (plotted on the x-axis in the following graph). Our simulations end when all soft tasks have been served.

We experiment on a GPL home made event-based simulator.

### 5.2 Results

Our simulations show an improvement over a simple BS. However, we are nearer from a BS than from the same policy using an exact computation of the slack. Moreover, the higher the periodic charge is, the worst the results are.

## 6. CONCLUSIONS

In this paper, we address the problem of jointly scheduling hard periodic tasks and soft aperiodic events with RTSJ. We present the programming constraints of RTSJ and propose a simple slack time evaluation algorithm. We describe a possible implementation of this algorithm. We describe a RTSJ compliant slack utilization algorithm. We simulate this algorithm with an exact slack computation and with our approximation. If we obtain better results than a simple BS, these simulations point out that our approximation of the available slack is too large.

## 7. REFERENCES

- [1] J. P. Lehoczky and S. Ramos-Thuel. An optimal algorithm for scheduling soft-aperiodic tasks fixed priority preemptive systems. In *proceedings of the 13th IEEE Real-Time Systems Symposium*, pages 110–123, Phoenix, Arizona, December 1992.
- [2] D. Masson and S. Midonnet. The design and implementation of real-time event-based applications with RTSJ. In *WPDRTS (in proceedings of IPDPS)*, page 148, Long Beach, CA USA, March 2007.